

Matching “versus” Mechanism Design

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Mechanism design approach

- ▶ Max objective s.t. constraints (technology, incentives)
- ▶ Vickrey auction
- ▶ Myerson auction

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Matching approach

- ▶ Seek a mechanism that satisfies "good properties"
- ▶ Gale-Shapley deferred acceptance algorithm
- ▶ Gale's Top Trading Cycles algorithm

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 - ▶ unique mechanism that is Pareto efficient and strategyproof
- ▶ Deferred acceptance can be formulated as a constrained optimization problem
 - ▶ Maximize proposer-side welfare s.t. stability constraints
 - ▶ G-S have a section in their paper on "optimality" that explicitly makes this point

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But, of course, often times these approaches end up looking quite different. (Else, Rakesh wouldn't have suggested this topic!)

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For all of these reasons it can be useful to find one or more good solutions to the problem, rather than pursue the optimal solution.

Keep in mind: Myerson, Vickrey ... these are the ones that worked!

- ▶ If only all problems had such elegant and compelling solutions.

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3. A few methodological observations
4. From theory to practice: course allocation at Wharton

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2. Gale Top Trading Cycles variant
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- ▶ “In other applications, the top trading cycles mechanism may be more appealing. . . .”
- ▶ “In other cases the choice between the two mechanisms may be less clear and it depends on the policy priorities of the policy makers”

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Moreover, subsequent empirical work on NYC suggests that the good properties approach was a reliable guide to welfare

The fact that we don't know the "optimal" school choice mechanism doesn't mean that we shouldn't discuss "good" school choice mechanisms!

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Other examples: assigning interchangeable workers to tasks or shifts; leads to salespeople; takeoff and landing slots to airlines; shared scientific resources amongst scientists; players to teams

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 - ▶ Will sometimes make additional assumptions about preferences (e.g. responsiveness)
- ▶ An allocation $\mathbf{x} = (x_i)_{i=1}^N$ is feasible if each $x_i \in 2^{\mathcal{C}}$ and $\sum_{i=1}^N x_{ij} \leq q_j$ for each j

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By contrast, in TU settings the three concepts tend to exactly coincide (e.g. Vickrey auction)

Impossibility Theorems

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 - ▶ Note contrast to setting with monetary transfers; VCG maximizes social welfare and is strategyproof

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- ▶ Essentially no progress on the "constrained Max SWF" problem, for either Bayesian IC or dominant strategy IC

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- ▶ Hatfield (2009, p. 514): "[the] results have shown that the only acceptable mechanisms for allocation problems of this sort is a sequential dictatorship, even when we restrict preferences to be responsive (...). Although unfortunate, it seems that in many of these applications, the best procedure (...) may well be a random serial dictatorship."

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- ▶ Both students agree that any "good" class is better than any "bad" class, and have responsive preferences
- ▶ Among the many ex-post Pareto efficient allocations are those in which one student gets all 10 good courses, while the other gets all 10 bad courses.

A mechanism from practice: the “draft”

Budish and Cantillon: "The Multi-Unit Assignment Problem: Theory and Evidence from Course Allocation at Harvard" (*AER*, 2012)

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 1. Students submit preferences, in the form of an ROL over courses (implicit assumption: preferences are responsive)
 2. Students are randomly ordered by the computer
 3. Students are allocated courses one at a time, based on their reported preferences and remaining availability.
 - ▶ Rounds 1, 3, 5, ...: ascending priority order
 - ▶ Rounds 2, 4, 6, ...: descending priority order

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Data (from 2005-2006 academic year)

- ▶ Students' actual submitted ROLs (potentially strategic)
- ▶ Students' underlying truthful ROLs, from an administration survey (caveats / robustness in paper)

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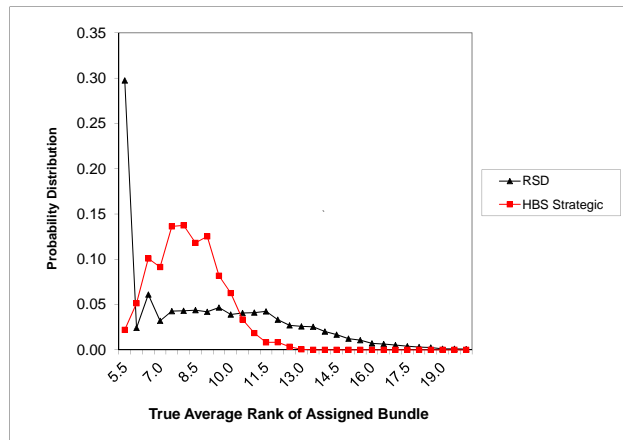
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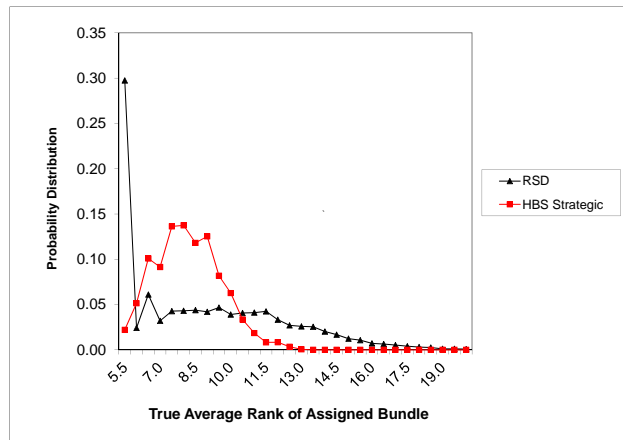
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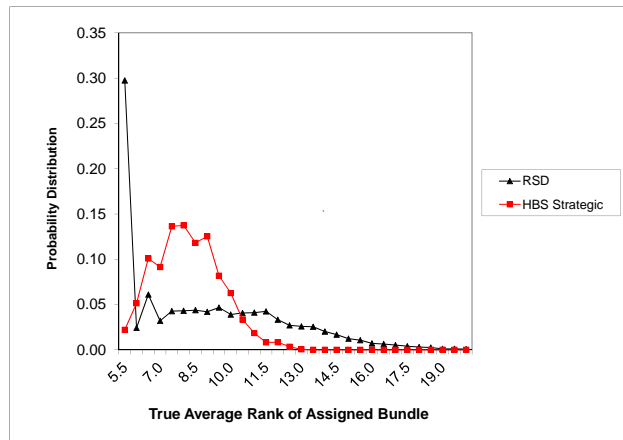
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- ▶ HBS Second-Order Stochastically Dominates RSD
- ▶ Implication: social planner prefers HBS to RSD if students have average-rank preferences and are weakly risk-averse

Why is RSD so unattractive ex-ante? Example

Suppose there are 4 courses with capacity of $\frac{1}{2}N$ seats each. Students require 2 courses each. Preferences are as follows:

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- ▶ Important note: unattractiveness of RSD does not depend on risk preferences. Even risk-neutral agents regard a "win a little, lose a lot" lottery as unappealing.

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- ▶ Overall, suggests a nuanced view of the role of strategyproofness in design, and the need for second-best alternatives to exact SP (eg “strategy-proofness in the large”, Azevedo and Budish, 2013)

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- ▶ Seek an incentives middle ground between strict strategyproofness (RSD) and simple-to-manipulate (HBS).
- ▶ Seek a mechanism that yields a relatively equal distribution of outcomes, like the draft and unlike the dictatorship

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No restrictions on preferences: students allowed to have arbitrary preferences over schedules. Allows for scheduling constraints, complementarities, etc.

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It is easy to see that existence is problematic.

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- ▶ Other extreme: dictatorships can be interpreted as exact CE, but from arbitrarily unequal budgets

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- ▶ Two agents. Four objects: two valuable Diamonds (Big, Small) and two ordinary Rocks (Pretty, Ugly). At most two objects per agent.
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 - ▶ Set other prices such that the poorer agent can afford {Small Diamond, Pretty Rock}, wealthier agent gets {Big Diamond, Ugly Rock}

Properties of the Approximate CEEI Mechanism

Efficiency

Ex-post efficient, but for small error

Fairness

$N+1$ Maximin Share Guarantee

Envy Bounded by a Single Good

Incentives

Strategyproof in the Large

Approximate CEEI and Matching “versus” Mechanism Design

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Approximate CEEI is attractive relative to alternatives under either interpretation. (As I stated at the outset, these approaches aren't always so different!)

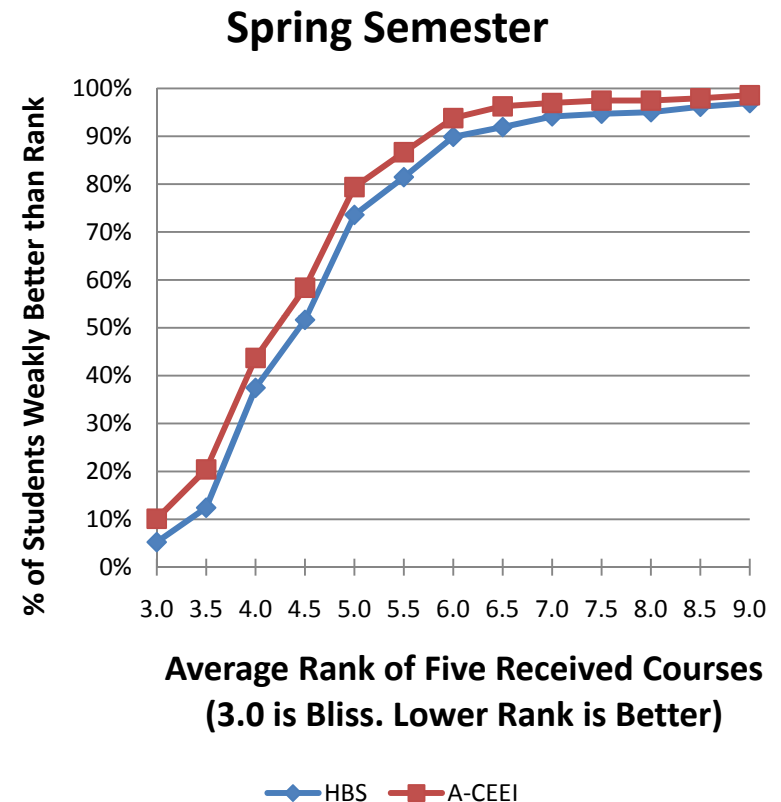
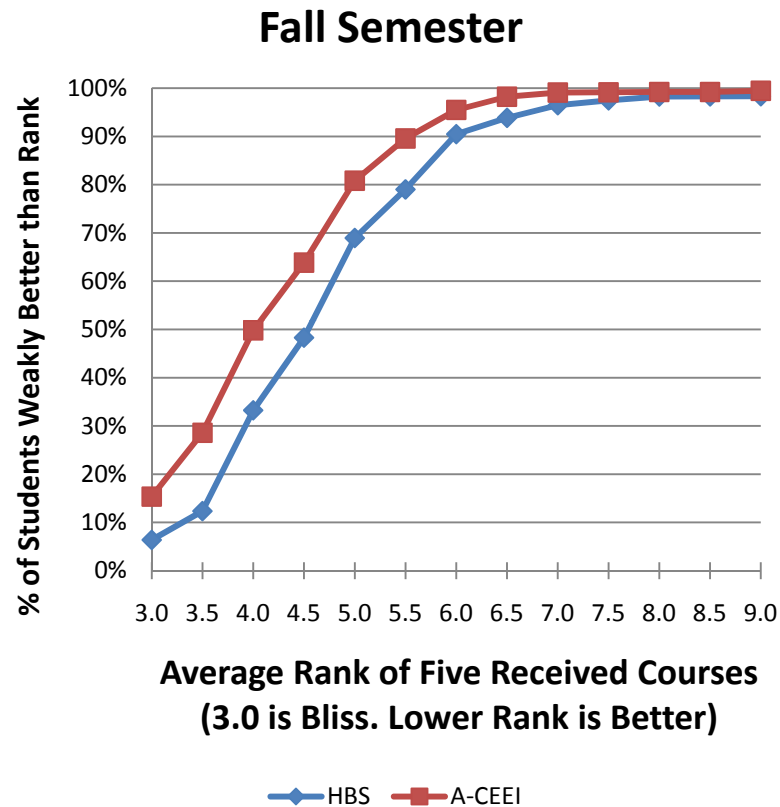
Table 2: Comparison of Alternative Mechanisms

Mechanism	Efficiency (Truthful Play)	Outcome Fairness (Truthful Play)	Procedural Fairness	Incentives	Preference Language
Approximate CEEI Mechanism (A-CEEI)	Pareto Efficient w/r/t Allocated Goods Allocation error is small for practice and goes to zero in the limit	N+1 – Maximin Share Guaranteed Envy Bounded by a Single Good	Symmetric	Strategyproof in the Large	Ordinal over Schedules
A-CEEI v2: Competitive Equilibrium from Equal-as-Possible Incomes (Sec 6.1)	Pareto Efficient	Worst Case: coincides with dictatorship	Symmetric	Strategyproof in the Large	Ordinal over Schedules
A-CEEI v3: A-CEEI with a Pareto- Improving Secondary Market (Sec 6.1)	Pareto Efficient	A bit weaker than N+1 – Maximin Share Guarantee, because prices in the initial allocation may be outside of $P(\delta, b')$. Initial allocation is Envy Bounded by a Single Good. The Pareto-improvement stage may exacerbate envy.	Symmetric	Manipulable in the Large	Ordinal over Schedules
Random Serial Dictatorship (Sec 8.1)	Pareto Efficient	Worst Case: Get k worst Objects	Symmetric	Strategyproof	Ordinal over Schedules
Multi-unit generalization of Hylland Zeckhauser Mechanism (Sec 8.2)	If vNM preferences are described by assignment messages, ex-ante Pareto efficient	If preferences are additive separable, envy bounded by the value of two goods Worst Case: Get Zero Objects	Symmetric	If vNM preferences are described by assignment messages, Strategyproof in the Large	Assignment messages
Bidding Points Mechanism (Sec 8.3)	If preferences are additive- separable, Pareto Efficient but for quota issues described in Unver and Sonmez (forth.)	Worst Case: Get Zero Objects	Symmetric	Manipulable in the Large	Cardinal over Items

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Sonmez-Unver (forth.) Enhancement to Bidding Points Mechanism	If preferences are additive-separable, Pareto Efficient	Worst Case: Get Zero Objects	Symmetric	Bidding Phase: Manipulable in the Large Allocation Phase: Strategyproof in the Large	Bidding Phase: Cardinal over Items Allocation Phase: Ordinal over Items
HBS Draft Mechanism (Sec 9.2)	If preferences are responsive, Pareto Efficient with respect to the reported information (i.e., Pareto Possible)	If preferences are responsive and $k=2$, Maximin Share Guaranteed If preferences are responsive, Envy Bounded by a Single Good	Symmetric	Manipulable in the Large	Ordinal over Items
Bezakova and Dani (2005) Maximin Utility Algorithm	If preferences are additive-separable, ideal fractional allocation is Pareto efficient. Realized integer allocation is close to the fractional ideal.	Worst Case: Get approximately zero objects (if a hedonist and all other agents are depressives)	Symmetric	Manipulable in the Large	Cardinal over items
Brams and Taylor (1996) Adjusted Winner	If preferences are additive-separable, Pareto Efficient	Worst Case: Get Zero Objects	Symmetric	Manipulable in the Large	Cardinal over Items
Herreiner and Puppe (2002) Descending Demand Procedure	Pareto Efficient	Does not satisfy Maximin Share Guarantee or Envy Bounded by a Single Object	Symmetric	Manipulable in the Large	Ordinal over Schedules
Lipton et al (2004) Fair Allocation Mechanism	Algorithm ignores efficiency	If preferences are additive separable, Envy Bounded by a Single Good	Symmetric	Manipulable in the Large	Cardinal over items
UChicago Primal-Dual Linear Programming Mechanism (Graves et al 1993)	Pareto Efficient when preference-reporting limits don't bind	Worst Case: Get Zero Objects	Symmetric	Manipulable in the Large	Cardinal over a Limited Number of Schedules

Figure 3: Ex-Ante Efficiency Comparison

Approximate CEEI Mechanism vs. HBS Draft Mechanism



Description: The Othman, Budish and Sandholm (2010) Approximate CEEI algorithm is run 100 times for each semester of the Harvard Business School course allocation data (456 students, ~50 courses, 5 courses per student). Each run uses randomly generated budgets. For each random budget ordering I also run the HBS Draft Mechanism, using the random budget order as the draft order. The HBS Draft Mechanism is run using students' actual strategic reports under that mechanism. The Approximate CEEI algorithm is run using students' truthful preferences. This table reports the cumulative distribution of outcomes, as measured by average rank, over the $456 \times 100 = 45,600$ student-trial pairs. Average rank is calculated based on the student's true preferences. For instance, a student who receives her 1,2,3,4 and 5th favorite courses has an average rank of $(1+2+3+4+5)/5 = 3$.

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- ▶ Perhaps there is a better way?
 - ▶ Recent work by Nguyen and Vohra seems quite promising in this regard
 - ▶ Instead of IC constraints, use envy-freeness constraints, and rely on large-market relationship between EF and SP-L (cf. Azevedo and Budish, 2013)

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 - ▶ We still don't know the “optimal” solution ...

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- ▶ We know from Micro 101 that we don't expect most preferences in the world to be lexicographic ... Perhaps we need new tools to make our preferences over mechanism designs a bit less lexicographic as well.

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- ▶ Comparing the good but imperfect mechanism to existing mechanisms gives a sense of the magnitude of the achievement
- ▶ Comparing the good but imperfect mechanism to a performance upper bound gives a sense of how much work there is left to do!